

Yices-QS 2022, an extension of Yices for quantified satisfiability

Stéphane Graham-Lengrand

SRI International, USA

1 Introduction

Yices-QS is a solver derived from Yices 2, an open-source SMT solver developed and distributed by SRI International. It extends Yices 2 to supports quantifiers for complete theories, and is unrelated to the support of quantifiers in Yices 2 for UF. Its core algorithm is a generalization of counterexample-guided quantifier instantiation (CEGQI) [Dut15] that can be seen as a form of lazy quantifier elimination. Yices-QS submits quantifier-free queries to Yices 2, leveraging some unique features pertaining to Yices’ MCSAT solver [dMJ13, Jov17]. Yices-QS is written in OCaml and uses the OCaml bindings for the Yices 2 C API:

<https://github.com/disteph/yicesQS>

https://github.com/SRI-CSL/yices2_ocaml_bindings

<https://github.com/SRI-CSL/yices2>.

In the 2022 SMT competition, Yices-QS entered logics BV, NRA, NIA, LRA, and LIA (single-query, non-incremental tracks), extending the scope of the 2021 submission (BV and NRA). The 2022 version is commit aa67ec2 of Yices-QS, with commit 989b4ad of the OCaml bindings and commit 09f1621 of Yices 2.

2 Algorithm

Yices-QS does not modify the structure of quantifiers in formulas: it does not prenexify formulas and, more importantly, it does not skolemize them to avoid introducing uninterpreted function symbols. In that, Yices-QS departs from the standard architecture for quantifier support consisting of keeping a set of universally quantified clauses, to be grounded and sent to a core SMT-solver for ground clauses. Instead, it sees a formula as a 2-player game, in the tradition of Bjørner & Janota’s *Playing with Quantified Satisfaction* [BJ15] and earlier work on QBF. Yices-QS’s algorithm is designed to answer queries of the following form:

“Given a formula $A(\bar{z}, \bar{x})$ and a model $\mathfrak{M}_{\bar{z}}$ on \bar{z} , produce either

- SAT($U(\bar{z})$), with $U(\bar{z})$ under-approx. of $\exists \bar{x} A(\bar{z}, \bar{x})$ satisfied by $\mathfrak{M}_{\bar{z}}$; or
 - UNSAT($O(\bar{z})$), with $O(\bar{z})$ over-approx. of $\exists \bar{x} A(\bar{z}, \bar{x})$ not satisfied by $\mathfrak{M}_{\bar{z}}$;
- where under-approximations and over-approximations are quantifier-free.”

To answer such queries, Yices-QS calls Yices’s feature *satisfiability modulo a model*, while the production of under- and over-approximations leverages *model interpolation* and *model generalization*.

When the input formula is in the exists-forall fragment, the algorithm degenerates to the one used in Yices’ $\exists\forall$ solver, using quantifier-free solving and *model generalization*, as described in [Dut15]. *Model interpolation*, a form of which is used within MCSAT to solve quantifier-free problems, also becomes useful with more quantifier alternations than $\exists\forall$. It generalizes to non-Boolean inputs the notion of UNSAT cores, which has been used in the *quantified-problems-as-games* approach [BJ15].

3 Theory-specific aspects

- *Model interpolation* is available in Yices’s MCSAT solver for quantifier-free formulas. In particular, it has theory-specific techniques for arithmetic, based on NLSAT [JdM12] (and ultimately, Cylindrical Algebraic Decomposition–CAD), and bitvectors [GLJD20].
- *Model generalization* for quantifier-free formulas can be done generically by substitutions [Dut15], but this can be complemented by theory-specific techniques that can provide better generalizations. For arithmetic, we use model-projection (based on CAD once again), in combination with generalization-by-substitution. For bitvectors, we use invertibility conditions from Niemetz et al. [NPR⁺18], including ϵ -terms, in combination with generalization-by-substitution. For the BV theory, the cegqi solver from [NPR⁺18] is probably the closest to Yices-QS, which approaches BV as an instance of the theory-generic algorithm from Section 2.

Note: For NRA, the presence of division in benchmarks departs from the theoretic applicability of Yices-QS’s algorithm for complete theories, because of division-by-zero (which also makes the theory undecidable). Yices’s CAD-based model-projection in NRA does not support division. In practice, when Yices-QS needs to perform model generalization with a formula involving division, it cannot use CAD model-projection and resorts to generalization-by-substitution. Resorting to generalization-by-substitution for NRA also means that Yices-QS’s algorithm may not terminate.

2022 vs 2021: Beyond a few bug fixes and optimizations, the 2022 version differs from the 2021 one in that it supports algebraic numbers in the generalization-by-substitution approach. The 2022 version can substitute a variable by its algebraic value in the model (represented with an ϵ -term), while the 2021 version would fail and give up.

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