

# Bit-Blasting for SMT-NIA with AProVE\*

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In automated termination provers like our tool AProVE [4], often the need arises to solve Boolean combinations of constraints in non-linear (integer) arithmetic to perform a proof step for a successful termination proof. Examples for prominent termination techniques where this is the case are well-founded orders based on polynomial [3] or matrix interpretations [2]. In order to facilitate this task, AProVE features a dedicated SMT-solver for SMT-NIA. The approach we are using at SMT-COMP for the QF\_NIA category is based on a reduction to a SAT problem (see [3] for details) and a subsequent invocation of the SAT solver MiniSAT[1]. This kind of approach is commonly known as *bit-blasting*.

Since satisfiability of SMT-NIA is undecidable in general, we search for solutions from a finite domain for the variables in the SMT-NIA instances. In case a satisfying assignment is found by the SAT solver, we return the corresponding integer solution. If the SAT solver detects unsatisfiability of the generated SAT instance, we know that there is no solution for the SMT-NIA instance with the used domain, and we restart the search with an extended domain. In this way, we obtain a semi-decision procedure for SMT-NIA. We additionally use information from global constraints (i.e., atomic top-level assertions like  $x \geq 42$ ) of the SMT-NIA instance to bound the search space for certain variables.

## References

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\* System description for *SMT-COMP 2015*.